Week 10: Lecture B Network Denial of Service

Thursday, October 30, 2025

Announcements

- Project 3: WebSec released
 - Deadline: Thursday, November 6th by 11:59PM (next week)

Project 3: Web Security

Deadline: Thursday, November 6 by 11:59PM.

Before you start, review the course syllabus for the Lateness, Collaboration, and Ethical Use policies.

You may optionally work alone, or in teams of at most two and submit one project per team. If you have difficulties forming a team, post on Piazza's Search for Teammates forum. Note that the final exam will cover project material, so you and your partner should collaborate on each part.

The code and other answers your group submits must be entirely your own work, and you are bound by the University's Student Code. You may consult with other students about the conceptualization of the project and the meaning of the questions, but you may not look at any part of someone else's solution or collaborate with anyone outside your group. You may consult published references, provided that you appropriately cite them (e.g., in your code comments). **Don't risk your grade and degree by cheating!**

Complete your work in the **CS 4440 VM**—we will use this same environment for grading. You may not use any **external dependencies**. Use only default Python 3 libraries and/or modules we provide you.



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Project 3 progress

| Working on Part 1 | |
|------------------------------------|----|
| | 0% |
| | |
| Finished Part 1, working on Part 2 | |
| | 0% |
| | |
| Finished Part 2, working on Part 3 | |
| | 0% |
| | |
| Finished with everything! | |
| | 0% |
| | |
| Haven't started yet :(| |
| | 0% |
| | |



Last time on CS 4440...

Application Layer Attacks
HTTP Content Injection
SMTP Header Spoofing
DNS Hijacking

Application Layer Attacks

- Application Layer:
 - ???

Application Layer Attacks

- Application Layer: where network-facing apps send/receive message
 - Application-specific protocols (message semantics, structure, processing rules, etc.)
- Attacking the application layer:
 - ???





6

Application Layer Attacks

- Application Layer: where network-facing apps send/receive message
 - Application-specific protocols (message semantics, structure, processing rules, etc.)
- Attacking the application layer:
 - Command Injection
 - SQL injection, CSRF, XSS
 - Denial of Service
 - Crash a remote application
 - Prevent others from using it
 - Message Tampering / Sniffing
 - Injecting data into messages
 - Capturing unencrypted data
 - Other protocol-specific attacks





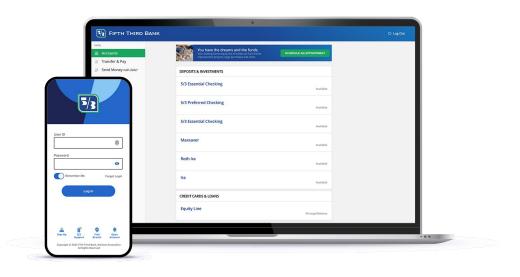




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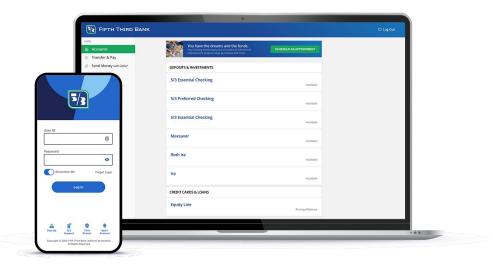
The HTTP Protocol

What is HTTP?



The HTTP Protocol

- What is HTTP?
 - Protocol for transmitting hypermedia documents (e.g., web pages)
- HTTP's Characteristics:
 - Widely used
 - Simple
 - Unencrypted



HTTP Tampering

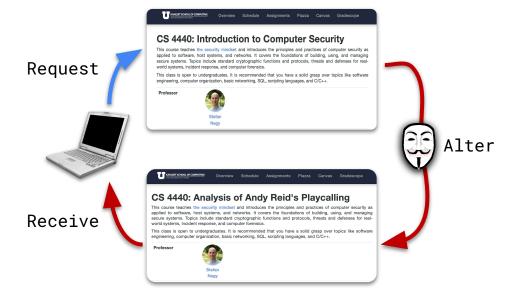
Attack: ???



Alter

HTTP Tampering

- Attack: exploit HTTP's insecurity
 - Nothing is encrypted!
- Attacker intercepts requested webpage and modifies it
 - User receives modified webpage
- Attacker capabilities?
 - Inject malicious content
 - Inject malicious code



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Thwarting HTTP Injection

Defenses: ???

Thwarting HTTP Injection

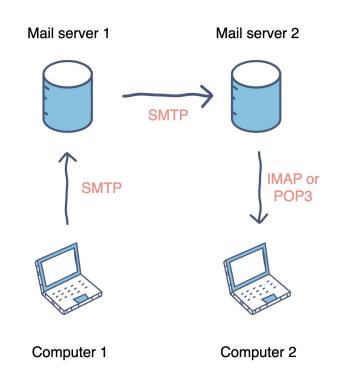
Defenses: encrypt everything all the time!

Answer: completely ditch HTTP!

- As web and app developers, enforce strict HTTPS compliance
 - Necessary to prevent HTTPS→HTTP downgrade attacks

13

- SMTP: Simple Mail Transfer Protocol
 - Implemented in the application layer
- Characteristics:
 - Text-based
 - Connection-oriented
 - Uses TCP ports 25/587
- Security guarantees:
 - ???



14

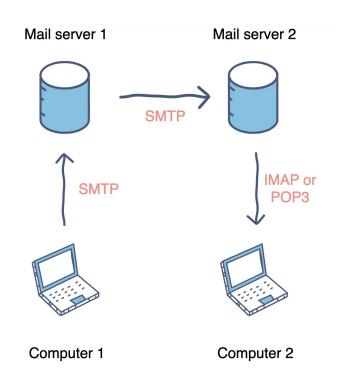
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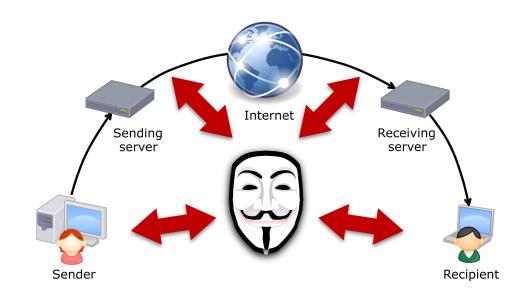
- Message integrity—no!
- Confidentiality—no!
- Authentication—no!



- No message integrity
 - ???

- No confidentiality
 - ???

- No authentication
 - ???



No message integrity

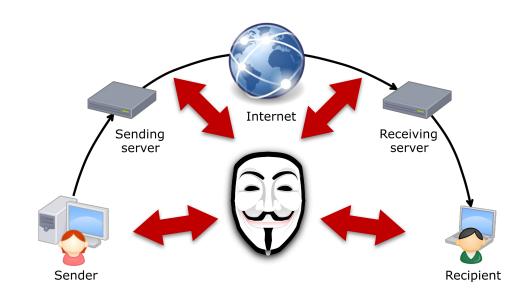
- Tamper with messages
- Block messages

No confidentiality

- Find sender/recipient
- Read message contents

No authentication

Spoof sender identity



SMTP Header Spoofing

Attack: ???

```
S: 220 attacker.com SMTP Exim
C: HELO attacker.com
S: 250 Hello attacker.com
C: MAIL FROM: <ceo@company.com>
S: 250 Ok
C: RCPT TO: <bob@company.com>
S: 250 Accepted
C: DATA
S: 354 Enter a message, ending with "." on a line by itself
C: Subject: Download this urgently
C: From: ceo@company.com
C: To: bob@company.com
C:
C: Hi Bob,
C: Please download this urgently: https://some-malicious-link.com
C: Regards
C: .
S: 250 OK
C: QUIT
S: 221 attacker.com closing connection
```

```
To: robert|bateman@email.com
Subject: Hi There
From: "Mickey Mouse" <m.mouse@disney.com>
X-Priority: 3 (Normal)
Importance: Normal
Errors-To: m.mouse@disney.com
Reply-To: m.mouse@disney.com
Content-Type: text/plain
```



18

SMTP Header Spoofing

Attack: spoof SMTP header to mislead recipient about sender of the email

```
S: 220 attacker.com SMTP Exim
C: HELO attacker.com
S: 250 Hello attacker.com
                                           Fake Sender
C: MAIL FROM: <ceo@company.com>
S: 250 Ok
                                           Victim
C: RCPT TO: <bob@company.com>
S: 250 Accepted
C: DATA
S: 354 Enter a message, ending with "." on a line by itself
C: Subject: Download this urgently
C: From: ceo@company.com
C: To: bob@company.com
C:
C: Hi Bob,
C: Please download this urgently https://some-malicious-link.com
C: Regards
C: .
                                   Malicious Link
S: 250 OK
C: OUIT
S: 221 attacker.com closing connection
```

```
To: robert bateman@email.com
Subject: Hi There
From: "Mickey Mouse" <m.mouse@disney.com>
X-Priority: 3 (Normal)
Importance: Normal
Errors-To: m.mouse@disney.com
Reply-To: m.mouse@disney.com
Content-Type: text/plain
```

Thwarting Email Spoofing

Defenses: ???

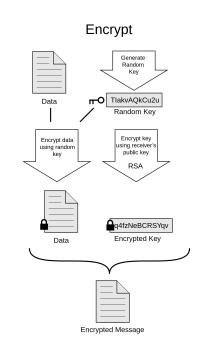
Thwarting Email Spoofing

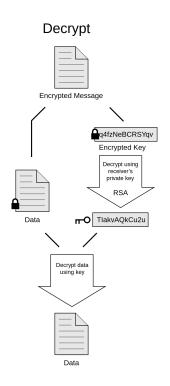
Checking email bodies

- Included links
- Attached files
- Text analysis (e.g., known spam campaigns)

Checking email headers

- Egress server domain registration
 - Check that sender is who it says it is
- Pretty Good Privacy (PGP)
 - Sender and Receiver authentication
 - Confidentiality
 - Integrity







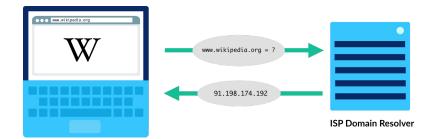
4

Identification on the Web

- How do we identify people?
 - Social security numbers
 - Passports, drivers licenses
 - Their unique fingerprints
- How can we identify internet hosts?
 - ???

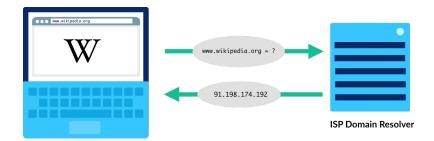
Identification on the Web

- How do we identify people?
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- How can we identify internet hosts?
 - Network layer: location via IP addresses
 - A logical addressing system
 - 32-bit (IPV4) addressing datagrams
 - What you care about: ???



Identification on the Web

- How do we identify people?
 - Social security numbers
 - Passports, drivers licenses
 - Their unique fingerprints
- How can we identify internet hosts?
 - Network layer: location via IP addresses
 - A logical addressing system
 - 32-bit (IPV4) addressing datagrams
 - What you care about: the domain name
 - E.g., www.wikipedia.org

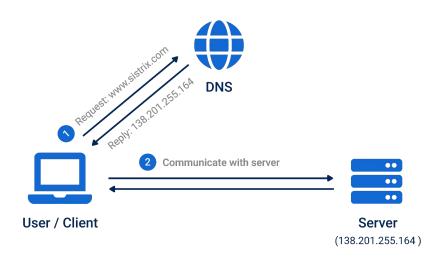




- Distributed database implemented in hierarchy of many name servers
- Application-layer protocol:
 - Hosts and domain name servers communicate to resolve domain names
 - Address-name translation
- Result: user requests ???
 - But their host really gets ???



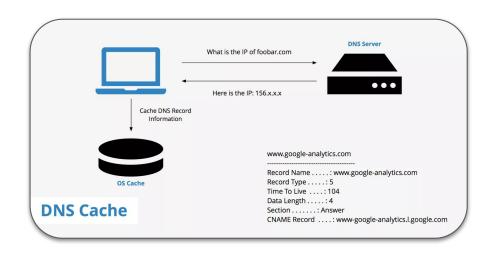
- Distributed database implemented in hierarchy of many name servers
- Application-layer protocol:
 - Hosts and domain name servers communicate to resolve domain names
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- Result: user requests domain name
 - But their host really gets its IP address
 - Convenient!



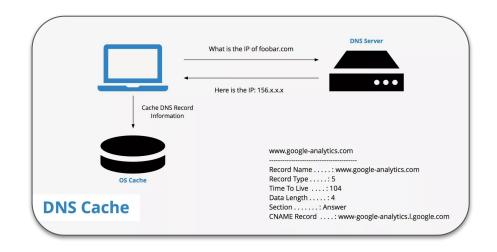
- How can we optimize DNS resolution?
 - ????

- How can we optimize DNS resolution?
 - Cache look-ups to amortize initial look-up, reduce system load
- Optimization 1: ???

Optimization 2: ???



- How can we optimize DNS resolution?
 - Cache look-ups to amortize initial look-up, reduce system load
- Optimization 1: temporal locality
 - www.espn.com/page1
 - www.espn.com/page2
- Optimization 2: popular domains
 - google.com
 - Facebook.com



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```
stefan@cs4440:~$ time nslookup facebook.com
                127.0.0.53
Server:
                127.0.0.53#53
Address:
Non-authoritative answer:
        facebook.com
Name:
Address: 31.13.70.36
        facebook.com
Name:
Address: 2a03:2880:f10d:83:face:b00c:0:25de
        0m0.474s
real
        0m0.000s
user
sys
        0m0.015s
```

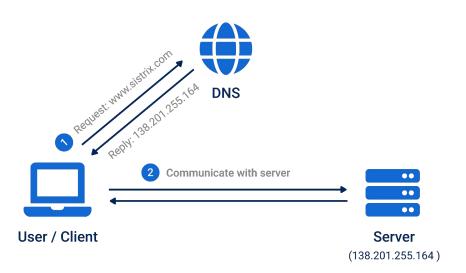
```
stefan@cs4440:~$ time nslookup facebook.com
                127.0.0.53
Server:
Address:
               127.0.0.53#53
Non-authoritative answer:
        facebook.com
Name:
Address: 31.13.70.36
        facebook.com
Name:
Address: 2a03:2880:f10d:83:face:b00c:0:25de
        0m0.023s
real
        0m0.000s
user
        0m0.011s
sys
```

First Lookup (non-cached)

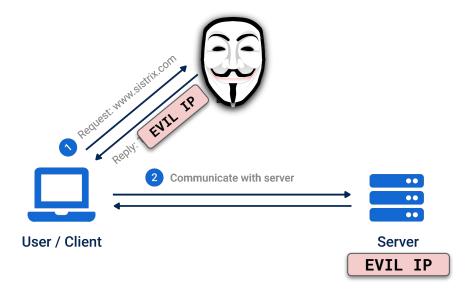
Second Lookup (cached)



What can an attacker do if they **control a DNS server**?



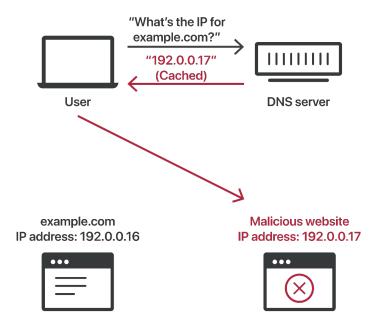
- What can an attacker do if they control a DNS server?
 - Control how users of that DNS server view the internet!
 - Assuming they use domain names



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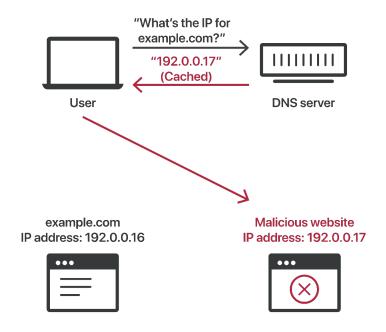
DNS Cache Poisoning

Attack: ???



DNS Cache Poisoning

- Attack: pre-empt DNS lookup by injecting malicious cache contents
 - Exploits DNS lookup optimization!
- Victim performs cache lookup, instead gets malicious domain IP
 - Attacker can redirect the victim's browser to the malicious website

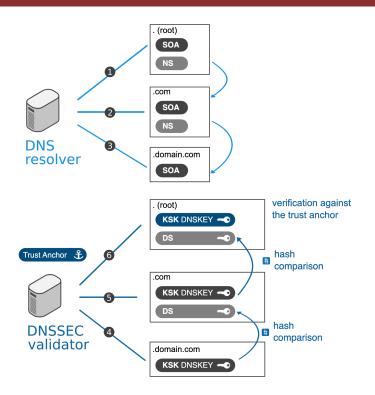


Thwarting DNS Hijacking

Defenses: ???

Thwarting DNS Hijacking

- DNS-level authentication
 - DNSSec
 - Public-key crypto to "sign" DNS records
- Endpoint authentication
 - Certify that what I am seeing really is bank.com
 - Transport Layer Security (TLS)



Questions?



This time on CS 4440...

Network Availability
Denial of Service (DoS) Attacks
Transport, Link, Network, and Physical DoS

Confidentiality: ???

Authenticity: ???

Integrity: ???

Access Control: ???

- Confidentiality: Concealment of information or resources
 - Attacks: ???
- Authenticity: Identification and assurance of info origin
 - Attacks: ???
- Integrity: Preventing improper and unauthorized changes
 - Attacks: ???
- Access Control: Enforce who is allowed access to what
 - Attacks: ???



- Confidentiality: Concealment of information or resources
 - Attacks: intercept credentials, info
- Authenticity: Identification and assurance of info origin
 - Attacks: SMTP header spoofing
- Integrity: Preventing improper and unauthorized changes
 - Attacks: tampering HTML over HTTP
- Access Control: Enforce who is allowed access to what
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 - Attacks: tampering HTML over HTTP
- Access Control: Enforce who is allowed access to what
 - Attacks: web app code injection
- Availability: Ability to use desired information or resource
 - Attacks: denial of service



Denial of Service Attacks

DoS: Denial of Service

Goal: ???



DoS: Denial of Service

Goal: make a service unusable, usually by overloading the server or network

How?

DoS: Denial of Service

- **Goal:** make a service unusable, usually by overloading the server or network
- How?
 - Trigger the host to crash
 - Application-based DoS
 - Memory corruption
 - Consume host's resources
 - TCP SYN floods
 - ICMP ECHO (ping) floods
 - Consume host's bandwidth
 - UDP floods
 - ICMP floods

Amazon loses **\$66,000 per minute** of downtime

Higher security makes DoS attacks **more likely**

Common DoS Attacks

- Locally-induced crash
 - Exploit host's OS or server software bug
- Local resource consumption
 - fork() bomb, fill disk, deep directory nesting
- Deny service to individual hosts
 - Force crash or outage of critical services
- Remotely-induced crash
 - "Magic" packets—ping of death, teardrop
- Remote resource consumption
 - Syslog, SYN, fragment flood, UDP storm





Common DoS Attacks (cont.)

- Deny service to an entire network
 - Target vulnerable links, critical network infrastructure
- Remotely-induced network outage
 - Attacks against routers, DNS servers
 - Redirected routes—forged routing information
- Remote network congestion
 - Remote control of compromised hosts ("zombies")
 - Allows for coordinated flooding
 - Distributed Denial of Service (DDoS)





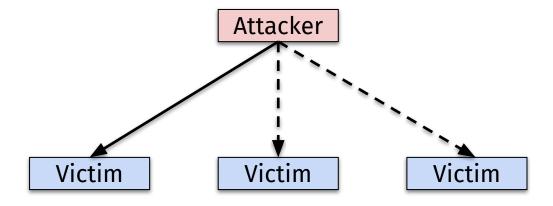
49

Simple DoS Attacks

Attacker spoofs their source address to hide origin

Defenses:

???

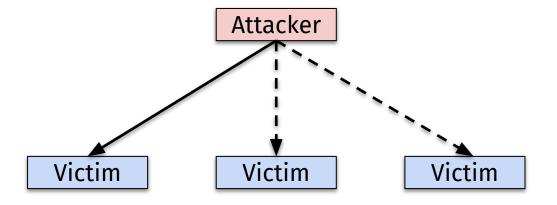


Simple DoS Attacks

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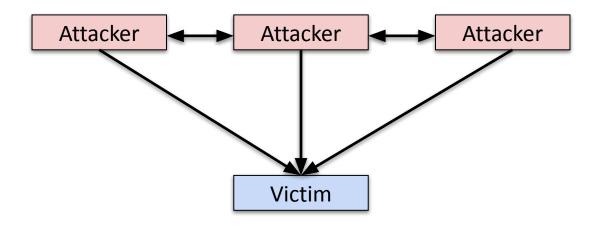
Defenses:

- Block source IP address
- Firewall
- ISP-level blocking
- Ignore requests

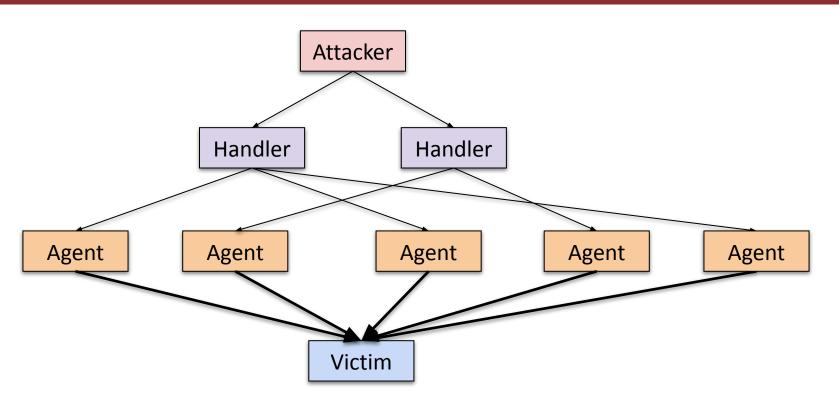


Coordinated DoS Attacks

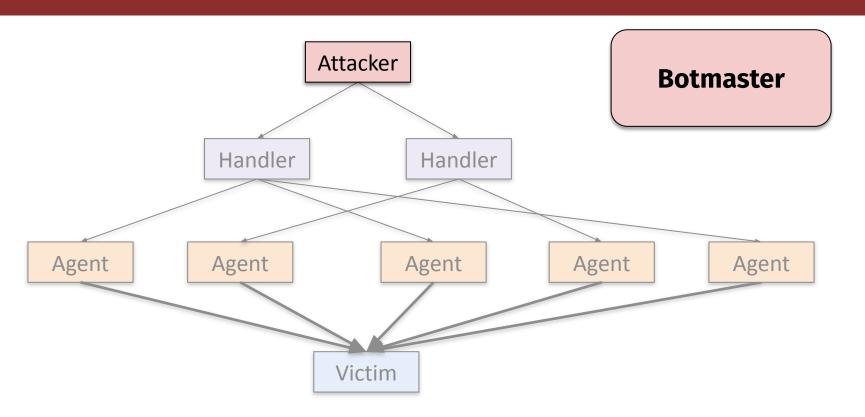
- Multiple willing attackers coordinate an attack on victim(s)
 - Same source-spoofing techniques as before
 - Harder to deal with



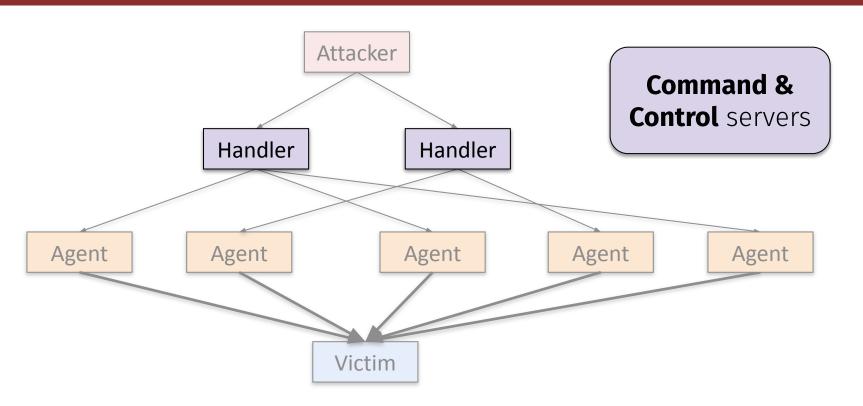
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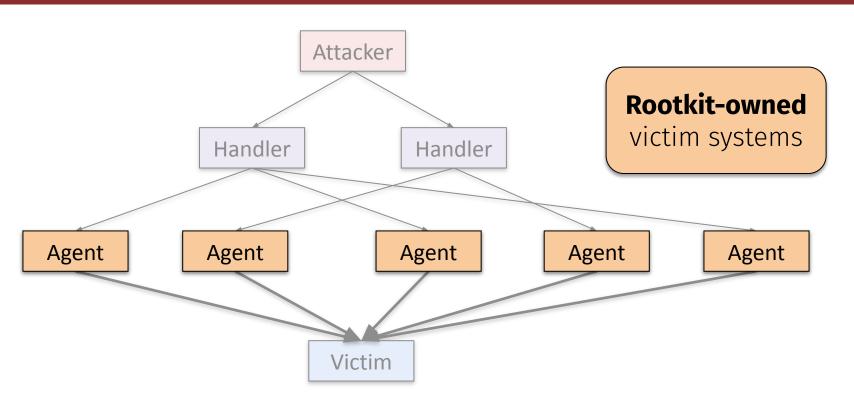














56

Timeline of a DDoS Attack

- Goal: compromise a large number of machines to form a botnet
 - 1. Attacker identifies **exploitable hosts** with scanners or other techniques
 - 2. Attacker gains control over systems via exploits, password cracking, etc.
 - Attacker installs rootkit
 - 4. Attacker remotely instructs compromised machines to attack the target

5

Timeline of a DDoS Attack

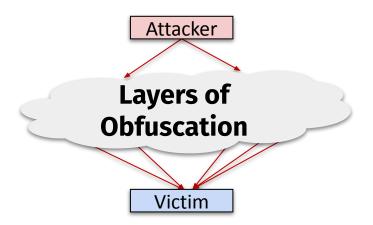
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Attacker



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Victim never sees the true attacker

Real DDoS Botnets

The Mirai botnet explained: How teen scammers and CCTV cameras almost brought down the internet

Mirai took advantage of insecure IoT devices in a simple but clever way. It scanned big blocks of the internet for open Telnet ports, then attempted to log in default passwords. In this way, it was able to amass a botnet army.

Storm: the largest botnet in the world?

Timely spam blasts help spread highly aggressive malware

A tiny botnet launched the largest DDoS attack on record

A small but powerful army of just 5,000 devices generated a record-breaking web attack.



60





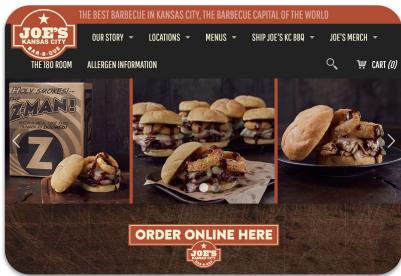


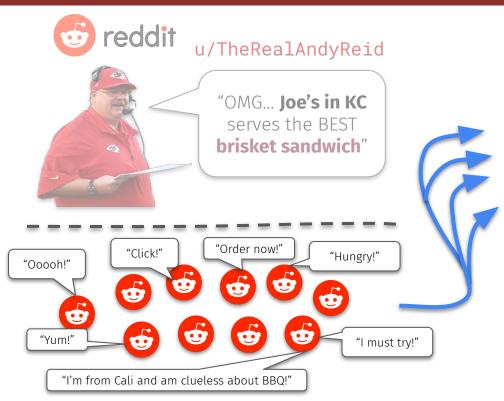
https://www.joeskc.com/





https://www.joeskc.com/





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How can we differentiate between Flash Mob traffic and DDoS traffic?

- How can we differentiate between Flash Mob traffic and DDoS traffic?
- Flash Mob traffic
 - Many clients using service legitimately
 - "Slashdot Effect", "Reddit Hug of Death"
 - Traffic dies down when the network is flooded
 - Sources in flash crowd are clustered
 - Usually by location (e.g., USA)

What Does Slashdot Effect Mean?

The slashdot effect refers to a temporary surge in traffic to a website, which can occur when a high-traffic website posts a link to smaller site or blog, thus directing an unprecedented surge in traffic. If the traffic increase is very large, it slow the site down or make it unreachable. The site is then considered to have been "slashdotted."

It's when someone posts a link to a website saying "Everyone, look at this website!" and everyone *does*. This puts so much traffic on the site in question's servers that they get overloaded and crash, causing the site to be inaccessible until the amount of traffic dies down a bit.



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How can we differentiate between Flash Mob traffic and DDoS traffic?

Flash Mob traffic

- Many clients using service legitimately
- "Slashdot Effect", "Reddit Hug of Death"
 - Traffic dies down when the network is flooded
- Sources in flash crowd are clustered
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DDoS traffic

- Attack does not end when host crashes
- Scattered locations (e.g., entire world)

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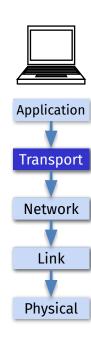
Questions?



Transport Layer DoS

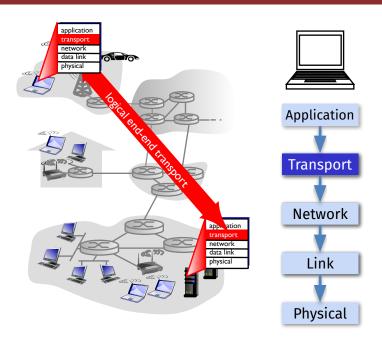
Recap: The Transport Layer

- What does it facilitate?
 - ???
- Key protocols?
 - Protocol 1: ???
 - Characteristics: ???
 - Protocol 2: ????
 - Characteristics: ???



Recap: The Transport Layer

- What does it facilitate?
 - Communication between apps on different hosts
- Key protocols?
 - Protocol 1: TCP (Transmission Control Protocol)
 - Characteristics: slow/complex but reliable
 - Protocol 2: UDP (User Datagram Protocol)
 - Characteristics: fast/simple but unreliable



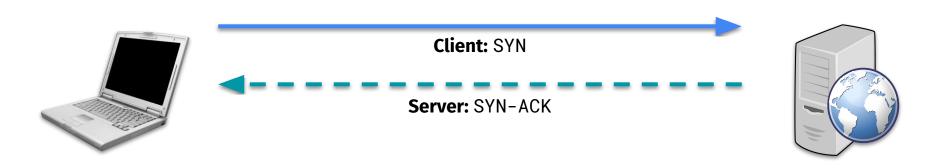
- Recall: TCP is a connection-oriented protocol
 - Initiate with three-way "handshake": SYN, SYN-ACK, ACK



Client: SYN

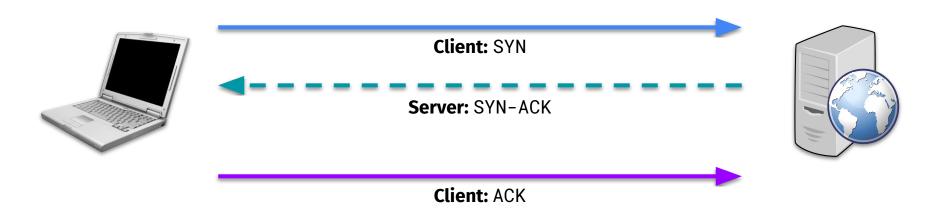


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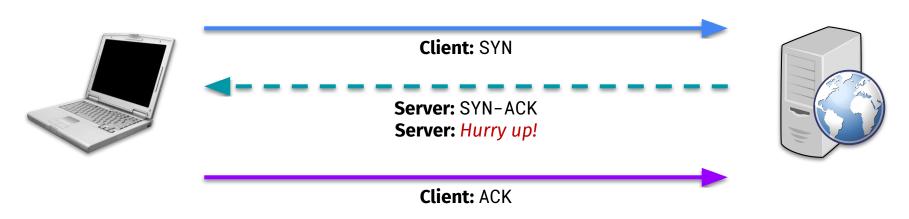


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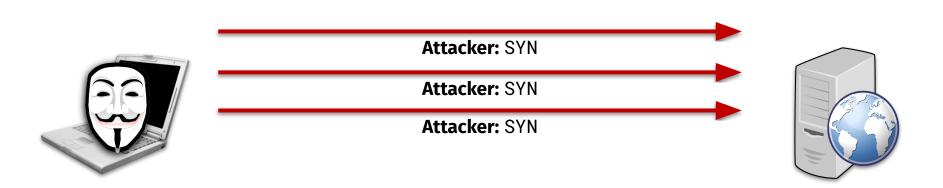
- Recall: TCP is a connection-oriented protocol
 - Initiate with three-way "handshake": SYN, SYN-ACK, ACK
 - Server waits until client responds with ACK





SYN Flooding Attack

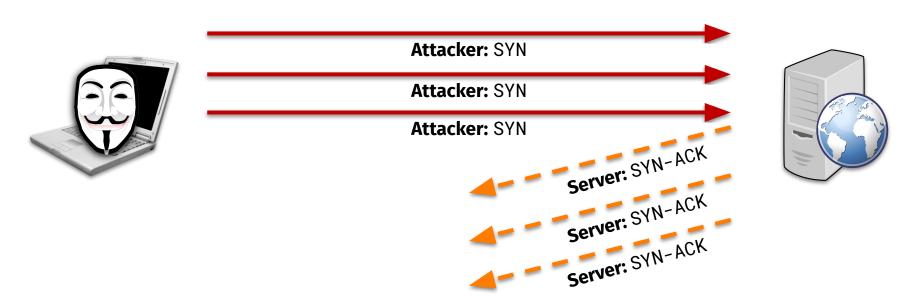
Attack: spam SYN packets to server, with spoofed origin address





SYN Flooding Attack

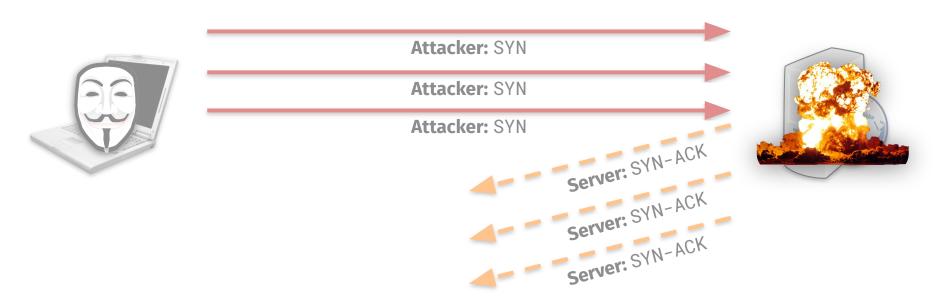
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SYN Flooding Attack

- Attack: spam SYN packets to server, with spoofed origin address
 - Server's resources completely reserved—now can't serve legitimate clients



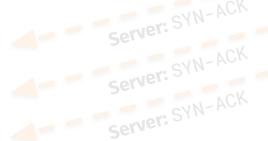
80

Thwarting SYN Flooding



How can we **prevent** SYN flooding?







Thwarting SYN Flooding

- Attack: spam SYN packets to server, with spoofed origin address
 - Server's resources completely reserved—now can't serve legitimate clients



How can we **prevent** SYN flooding?

Incorporate **state**—use **SYN cookies!**





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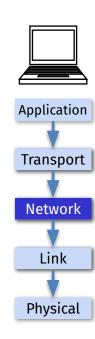
Questions?



Network Layer DoS

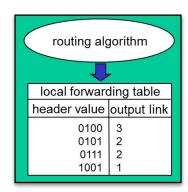
Recap: The Network Layer

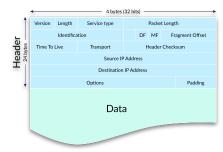
- What does it facilitate?
 - ???
- Key functions?
 - Function1: ???
 - Function1: ???
- Addressing?
 - ????

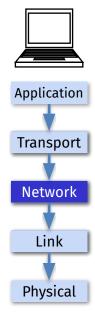


Recap: The Network Layer

- What does it facilitate?
 - Sending of data from host on one network to another
- Key functions?
 - **Function1:** Routing: (find the shortest path for a packet)
 - Function1: Forwarding (send packet on to the next hop)
- Addressing?
 - IP addressing (logical addressing)

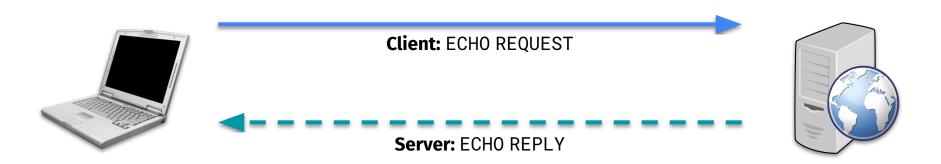






ICMP: Internet Control Message Protocol

- ICMP: pings to determine whether a system is connected to the Internet
 - Analogous to "Hello, are you still there?"



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ICMP Smurf Attacks

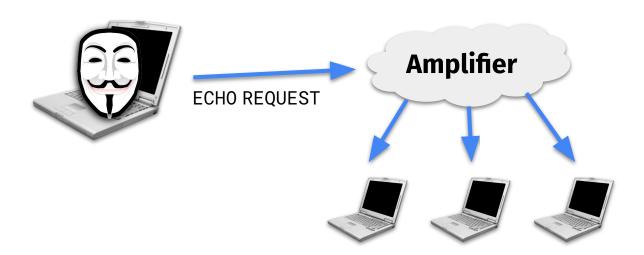
Attack: takes advantage of broadcast-enabled hosts to amplify attack





ICMP Smurf Attacks

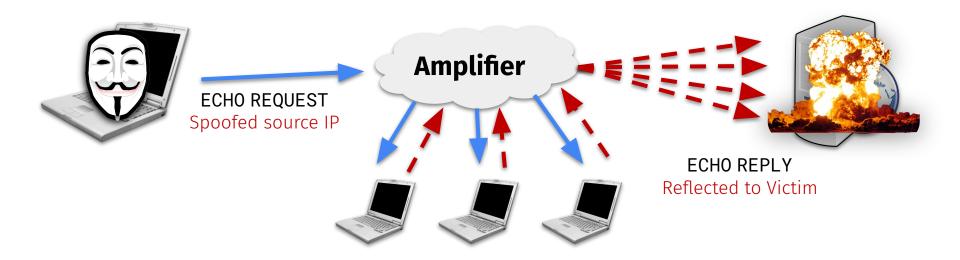
Attack: takes advantage of broadcast-enabled hosts to amplify attack





ICMP Smurf Attacks

- Attack: takes advantage of broadcast-enabled hosts to amplify attack
- Attacker spams spoofed-source ICMP requests, reflected to victim's IP



Advanced DoS Strategies

Reflection:

IP spoofing to redirect response to a victim

Amplification:

 Technique that increases the amount of traffic or packet size that the victim sees versus what the attacker originally sent

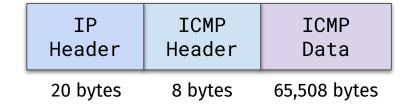
Common in real-world DDoS attacks

- Harder to detect (source obfuscation)
- Harder to thwart (changing sources)



ICMP Ping of Death Attack

- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in fragments and reassembled by receiver



ICMP Ping of Death Attack

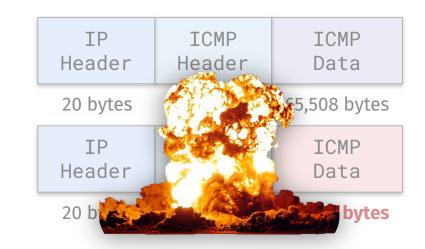
- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in fragments and reassembled by receiver
- Attack: send packet in fragments that reassemble to 64K+ bytes
 - Many historical computer systems
 could not handle larger packets

| 20 bytes | 8 bytes | 65,510 bytes |
|----------|---------|--------------|
| IP | ICMP | ICMP |
| Header | Header | Data |
| 20 bytes | 8 bytes | 65,508 bytes |
| IP | ICMP | ICMP |
| Header | Header | Data |

9.

ICMP Ping of Death Attack

- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in fragments and reassembled by receiver
- Attack: send packet in fragments that reassemble to 64K+ bytes
 - Many historical computer systems could not handle larger packets
- Result: crash by buffer overflow
 - Can't serve clients until restart!



Thwarting ICMP-based DoS

- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in fragments and reassembled by receiver
- How can we **prevent** ICMP attacks?
 - Many historical computer systems
 could not handle larger packets
- Result: crash by buffer overflow
 - Can't serve clients until restart!



Thwarting ICMP-based DoS

- Internet Protocol: IPV4 packets should be less than 65,536 bytes
 - Packets can be sent in fragments and reassembled by receiver
- Attack: sen that reasse

How can we **prevent** ICMP attacks?

Many his could not

Secure any open **amplifiers**, and **sanitize** network **input**

- Result: cras
 - Can't sei

Questions?



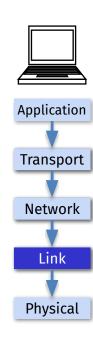
Link Layer DoS

98

Recap: The Data Link Layer

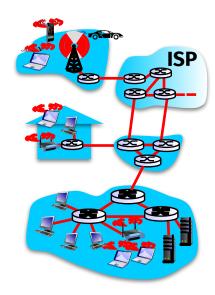
- What does it facilitate?
 - ????
- Addressing?
 - ????

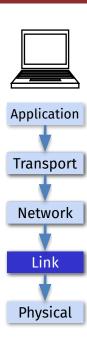
- Authenticity?
 - ???



Recap: The Data Link Layer

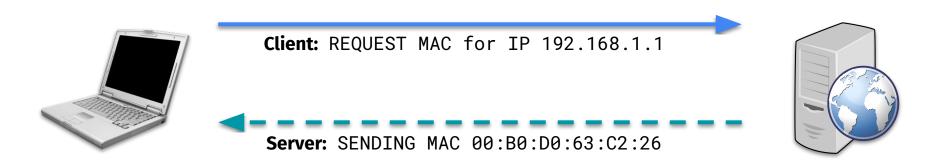
- What does it facilitate?
 - Responsible for the node-to-node delivery of data
- Addressing?
 - MAC addresses
 - Physical identifier for hardware
- Authenticity?
 - No—MAC addresses can be changed!





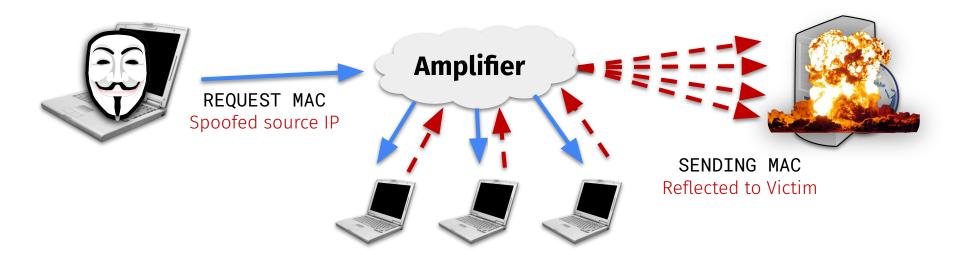
ARP: Address Resolution Protocol

- ARP: query to resolve the MAC address given a desired host IP
 - How we know which physical address to transmit data to from its logical address



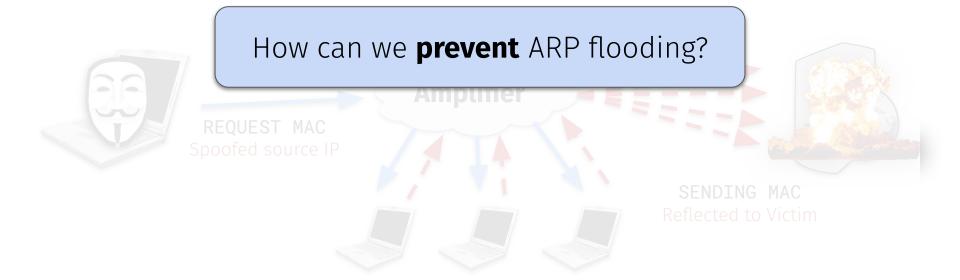
ARP Flooding Attack

- Attack: same idea as ICMP Smurfing; spoof source to victim and spam away!
 - Victim gets overwhelmed by ARP replies and bandwith crashes



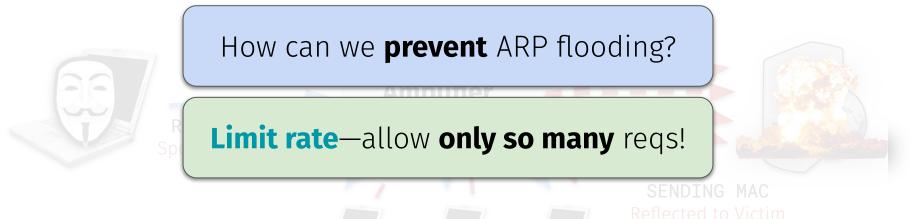
Thwarting ARP Flooding

- Attack: same idea as ICMP Smurfing; spoof source to victim and spam away!
 - Victim gets overwhelmed by ARP replies and bandwith crashes



Thwarting ARP Flooding

- Attack: same idea as SYN flood; spoof source to victim and spam away!
 - Victim gets overwhelmed by ARP replies and bandwith crashes





Physical Layer DoS

Recap: Physical Layer

What is it?

Recap: Physical Layer

What is it?











Provided by AT&T Wi-Fi



107

Physical Layer DoS

Russian Spy Submarines Are Tampering with Undersea Cables That Make the Internet Work. Should We Be Worried?

A massive cable attack is probably an over-hyped scenario, at least for a country with as many redundant cables as the United States pitted against a limited number of Russian special-operations submarines.



CNN Exclusive: FBI investigation determined Chinese-made Huawei equipment could disrupt US nuclear arsenal communications



Physical Layer DoS

Cables That Make the Internet

Work Iran blocks capital's internet access as Amini protests grow

> Social media platforms have also been cut off in areas of Tehran and Kurdistan as videos of dissent go viral

Chinese-made Huawei equipment could disrupt US nuclear arsenal communications

Thwarting Physical Layer Attacks



110

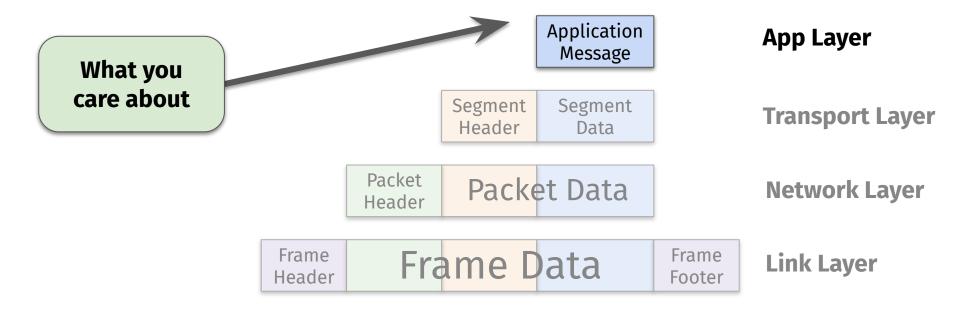
Questions?



Analyzing Network Packets

Recap: Internet Packet Encapsulation

How packets are generated and sent

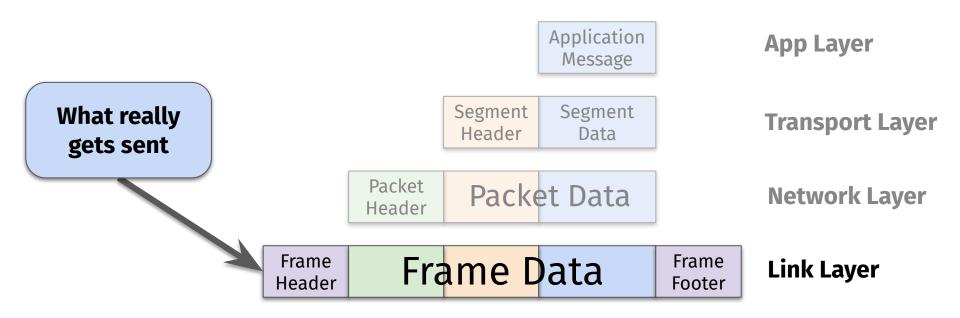




Stefan Nagy 113

Recap: Internet Packet Encapsulation

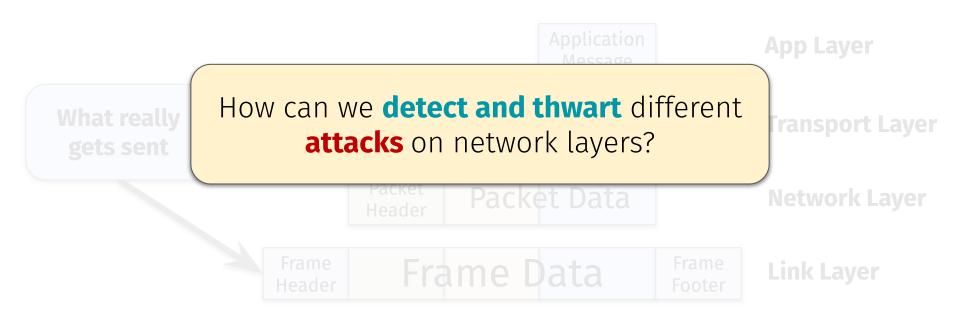
How packets are generated and sent





Recap: Internet Packet Encapsulation

How packets are generated and sent





Tools of the Trade

- Packet Analyzers:
 - Tools for dissecting network packets
- Packet Analyzers allow you to:
 - Identify unusual packets
 - Characterize network activity
 - Pinpoint malicious traffic
- The basis of modern-day network security (e.g., firewalls, antivirus)



Familiarity with packet analysis tools?

I eat NetSec CTF challenges like a kid eats candy on Halloween.

O%

Some (e.g., Wireshark, DPKT, Scapy, or something else)

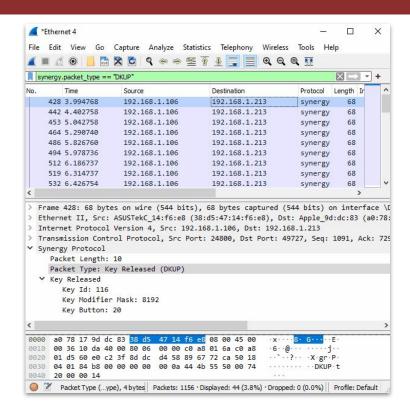
O%

None (but that's totally okay!)



Tools of the Trade: Wireshark

- A "graphical interface" for manual packet analysis
 - Completely open-source and free
- General workflow:
 - Load up a PCAP (packet capture)
 - Wireshark will display each packet
 - Inspect particular fields of interest



Tools of the Trade: Wireshark

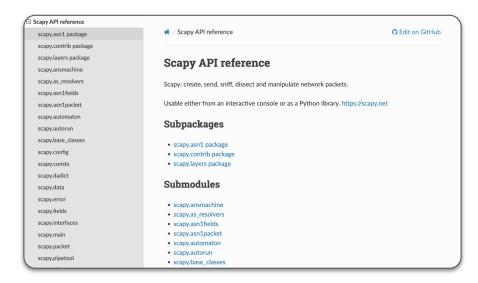
| No. | Time | Source | Destination | Protocol | Length Info |
|---|---------------------|---|---|----------|--|
| | 1 0.000000 | 10.0.0.2 | 10.128.0.2 | TCP | 54 3341 → 80 [SYN] Seq=0 Win=512 Len=0 |
| | 2 0.003987 | 10.128.0.2 | 10.0.0.2 | TCP | 58 80 → 3222 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 |
| | 3 0.005514 | 10.128.0.2 | 10.0.0.2 | TCP | 58 80 → 3341 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 |
| | 4 0.008429 | 10.0.0.2 | 10.128.0.2 | TCP | 54 3342 → 80 [SYN] Seq=0 Win=512 Len=0 |
| | 5 0.010233 | 10.128.0.2 | 10.0.0.2 | TCP | 58 80 → 3220 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 |
| | 6 0.014072 | 10.128.0.2 | 10.0.0.2 | TCP | 58 80 → 3342 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 |
| | 7 0.016830 | 10.0.0.2 | 10.128.0.2 | TCP | 54 3343 → 80 [SYN] Seq=0 Win=512 Len=0 |
| | 8 0.022220 | 10.128.0.2 | 10.0.0.2 | TCP | 58 80 → 3343 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 |
| | 9 0.023496 | 10.128.0.2 | 10.0.0.2 | TCP | 58 80 → 3219 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 |
| | 10 0.025243 | 10.0.0.2 | 10.128.0.2 | TCP | 54 3344 → 80 [SYN] Seq=0 Win=512 Len=0 |
| | 11 0.026672 | 10.128.0.2 | 10.0.0.2 | TCP | 58 80 → 3218 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 |
| | 12 0.028038 | 10.128.0.2 | 10.0.0.2 | TCP | 58 80 → 3221 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 |
| 41 | 13 0.030523 | 10.128.0.2 | 10.0.0.2 | TCP | 58 80 → 3344 [SYN, ACK] Seq=0 Ack=1 Win=29200 Len=0 MSS=1460 |
| 1 | cknowledgment n | 0 (relative s umber: 0 (rela umber: 1 (rela | equence number) tive sequence number)] tive ack number) | | |
| 0110 = Header Length: 24 bytes (6) | | | | | |
| Flags: 0x012 (SYN, ACK) | | | | | |
| Window size value: 29200 | | | | | |
| [Calculated window size: 29200] | | | | | |
| Checksum: 0x4268 [unverified] [Checksum Status: Unverified] | | | | | |
| | | | | | |
| | rgent pointer: | s), Maximum segme | nt size | | |
| | Timestamps] | s), maximum segme | III SIZE | | |
| | 1 Tille 2 Callips] | | | | , , , , , , , , , , , , , , , , , , , |
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Stefan Nagy 119

Tools of the Trade: Scapy

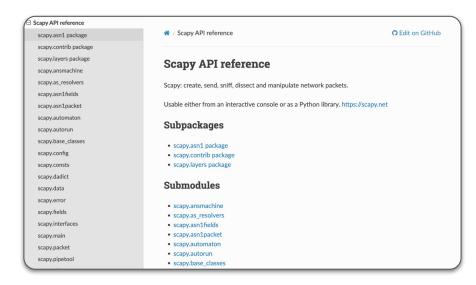
- Python API for programmatic packet capture and analysis
 - Think of it as "Wireshark in API form"
 - Project 4: you will use Scapy to write your own packet analysis scripts



120

Tools of the Trade: Scapy

- Python API for programmatic packet capture and analysis
 - Think of it as "Wireshark in API form"
 - Project 4: you will use Scapy to write your own packet analysis scripts
- We'll provide the PCAP traces...
 - You'll write code to analyze them!
 - Examples:
 - Detecting attacks on a network
 - Finding user credentials
 - Sniffing a user's browsing history



Questions?



Next time on CS 4440...

Passwords and Secure Authentication